

INTRO TO DYNAMIC PROGRAMMING



CPMSOC

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ATTENDANCE



MOTIVATING PROBLEM



The fibonacci numbers are defined $\text{fib}(0)=1$, $\text{fib}(1)=1$, and $\text{fib}(n)=\text{fib}(n-1)+\text{fib}(n-2)$.

e.g. 1,1,2,3,5,8,11,...

Given x , output the x -th fibonacci number.

RECURSION?

Implement the recursive
formula straight-forward.



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```
int fib(int n) {  
    if (n <= 1) return 1;  
    return fib(n - 1) + fib(n - 2);  
}
```

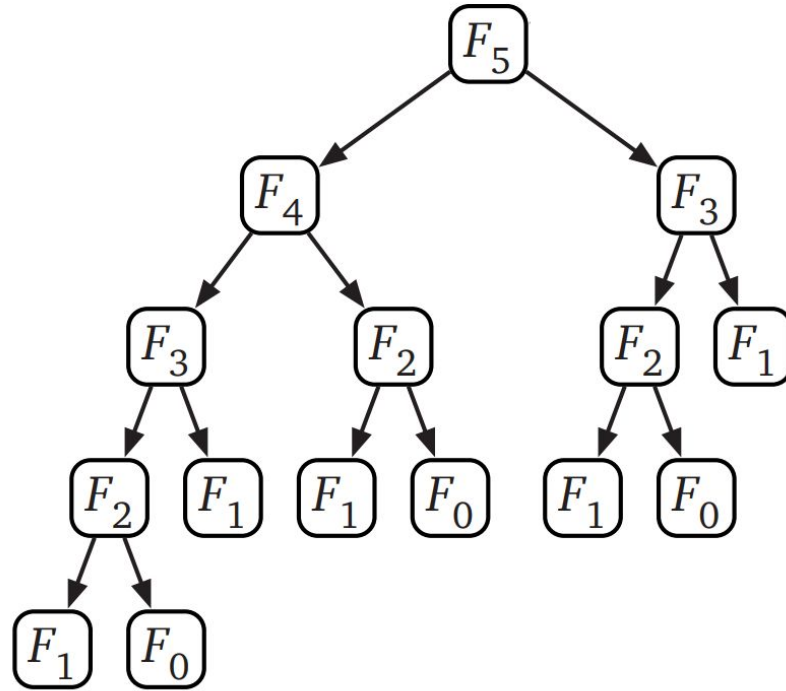
SO WHAT'S THE ISSUE?

Try to run for a number $x > 30$.

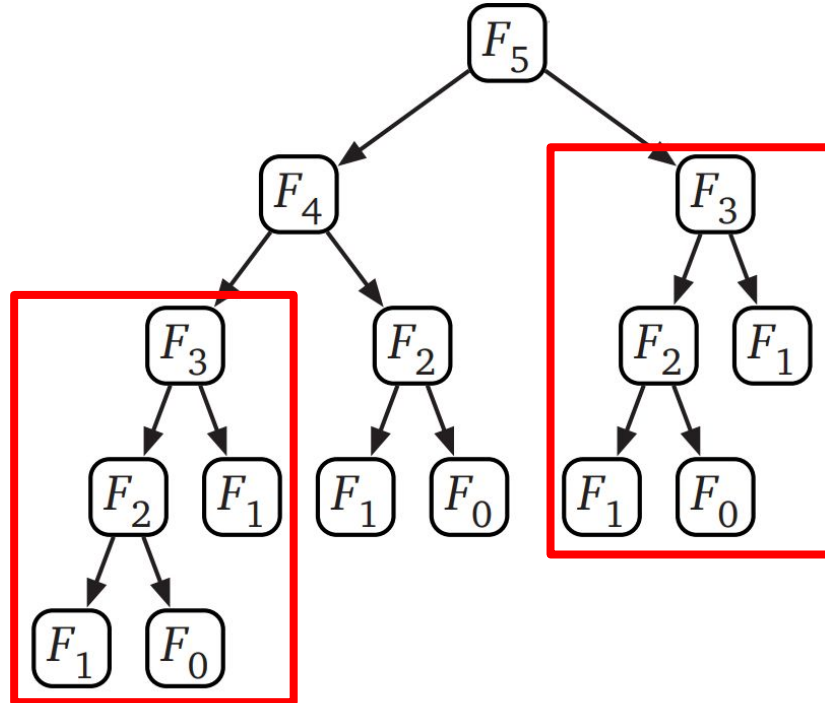
The program can't finish running...



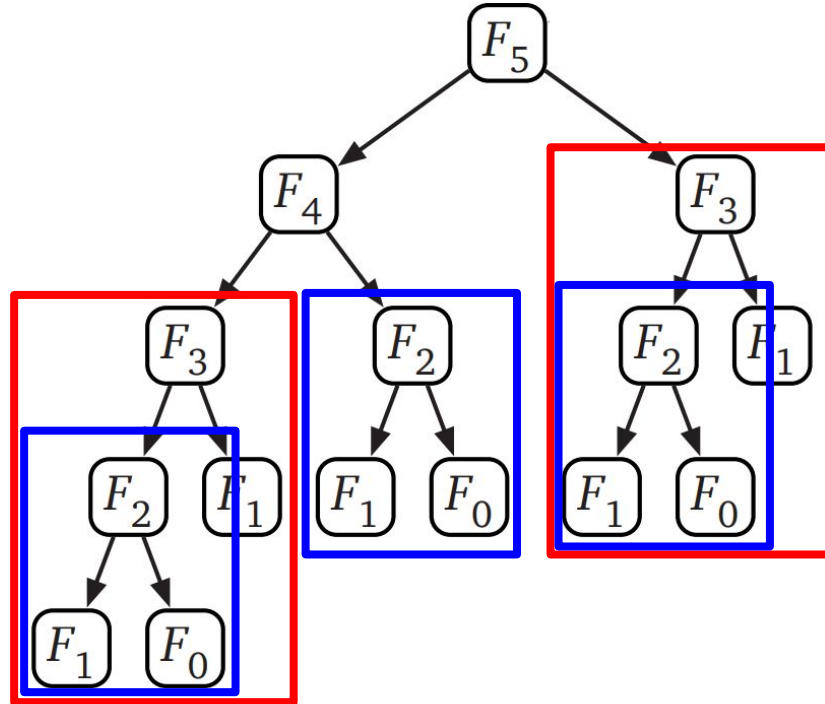
LET'S SEE WHAT'S HAPPENING



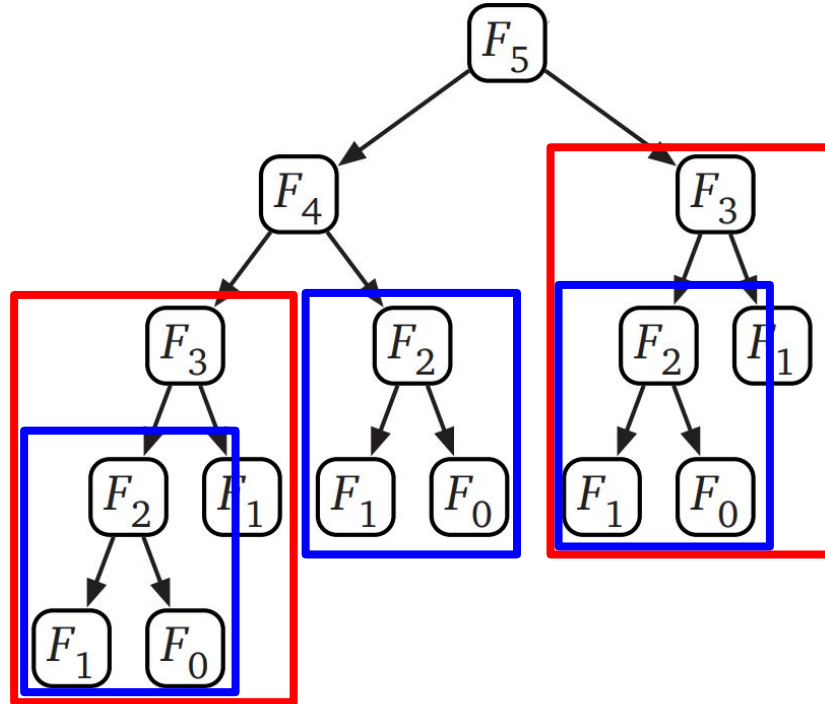
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Worst case we're calculating 2^x states!

$O(2^x)$

MEMOISATION!

Let's not recalculate the values we've already calculated before.

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```
int dp[MXX];  
  
bool seen[MXX];  
  
int fib(int n) {  
    if (n <= 1) return 1;  
    if (seen[n]) return dp[n];  
    dp[n] = fib(n - 1) + fib(n - 2);  
    return dp[n];  
}
```

WHY IS THIS BETTER?



We only calculate each state from $0 \sim x$ once. $O(x)$

This new method will now easily pass for much larger values of x . Yay!

WHAT INSIGHT DOES THIS GIVE TO DP?

- Similar sub-problems
- Similar sub-structures

We break the original problem up into smaller, manageable problems that we combine!



COIN CHANGE PROBLEM



Given some denomination of coins, and a value, what is the minimum number of coins to make that value?

Example:

Value: 12

Denominations: 1, 2, 5

IDEA



Let's always take the largest denomination coin that we can take, then move on to the next.

Example:

Value: 12

Denominations: 1, 2, 5

Solution: 3 coins (5, 5, 2)



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DOES THIS ALWAYS
WORK?

No

Consider:

Value: 11

Denominations: 1, 5, 7



No



Consider:

Value: 11

Denominations: 1, 5, 7

If we took greedily from biggest denomination: 5 coins (7, 1, 1, 1, 1)

Optimal solution: 3 coins (5, 5, 1)



ANY OTHER IDEAS?

LET'S THINK WITH DP



If we can't solve the problem for value V , let's solve it for value $V - c[i]$, where $c[i]$ is the denomination for each coin.

A recursive relation is being generated here!

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If we can't solve the problem for value V , let's solve it for value $V - c[i]$, where $c[i]$ is the denomination for each coin.

A recursive relation is being generated here!

$$dp[V] = \min\{dp[V - c[i]] + 1\}$$

WHEN DO WE STOP?

$dp[0] = 0$

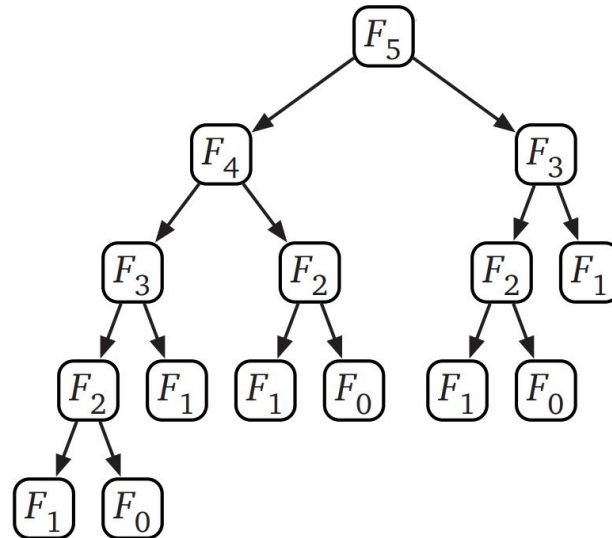
It requires 0 coins to form a total value of 0!



TOP DOWN DP



We recursively call “states” that may have not been calculated yet.



BOTTOM UP DP



We build our dp calculations upwards with the states that have already been calculated.

Let's see this in action with the Coin Change Problem!

COIN CHANGE PROBLEM VISUALISATION



Denominations: 1, 5, 7

0	1	2	3	4	5	6	7	8	9	10	11
0	1	2	3	INF	INF	INF	INF	INF	INF	INF	INF

COIN CHANGE PROBLEM VISUALISATION



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0	1	2	3	4	5	6	7	8	9	10	11
0	1	2	3	4	INF	INF	INF	INF	INF	INF	INF

COIN CHANGE PROBLEM VISUALISATION



Denominations: 1, 5, 7

0	1	2	3	4	5	6	7	8	9	10	11
0	1	2	3	4	1	INF	INF	INF	INF	INF	INF

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0	1	2	3	4	5	6	7	8	9	10	11
0	1	2	3	4	1	2	1	2	3	INF	INF

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Denominations: 1, 5, 7

0	1	2	3	4	5	6	7	8	9	10	11
0	1	2	3	4	1	2	1	2	3	2	INF

COIN CHANGE PROBLEM VISUALISATION



Denominations: 1, 5, 7

0	1	2	3	4	5	6	7	8	9	10	11
0	1	2	3	4	1	2	1	2	3	2	3

THANKS FOR ATTENDING!



Attendance ----->

